## SkyChain: A Deep Reinforcement Learning-Empowered Dynamic Blockchain Sharding System

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# Background





### **1** The bottleneck of scalability in blockchain

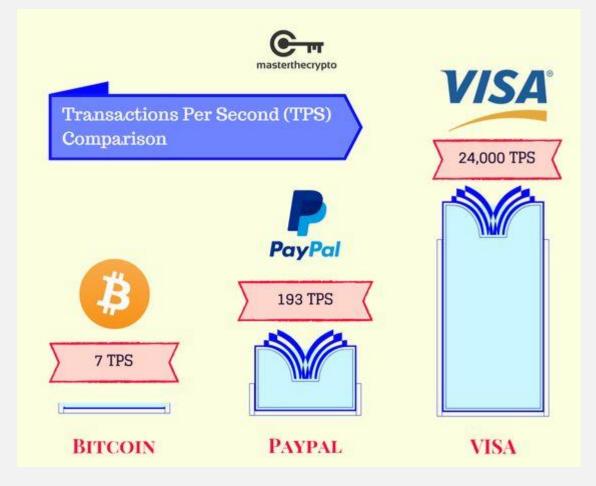
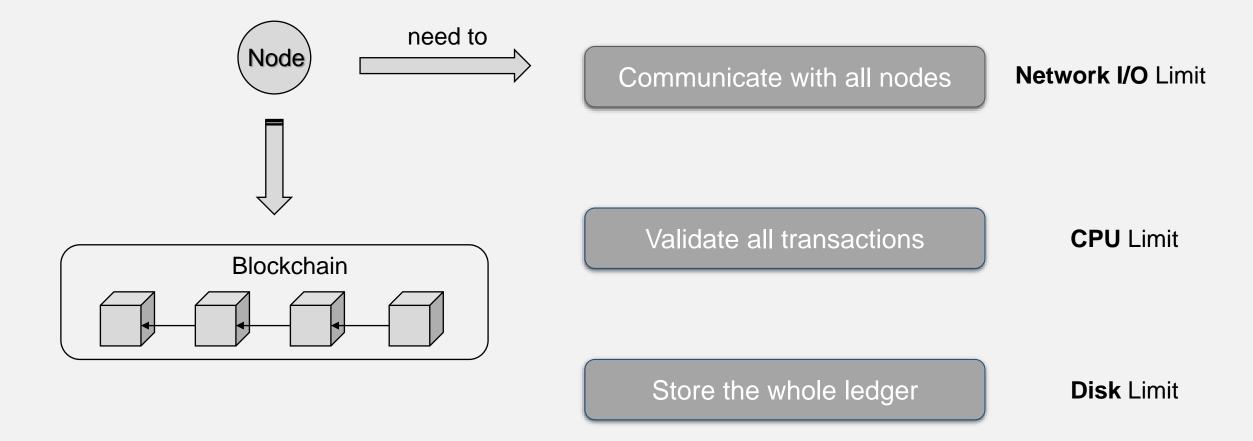


Fig.1:TPS Comparison

Low TPS

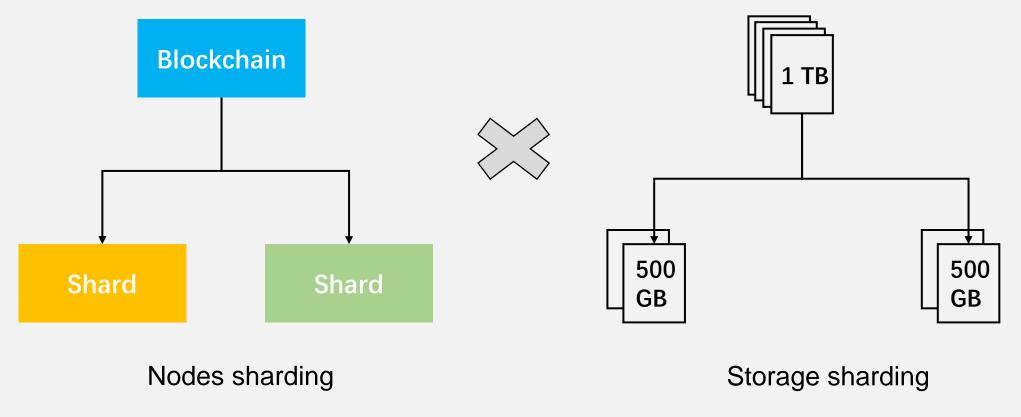


### 2 Why is it not scalable in blockchain?





### **3 Blockchain sharding overview**



Reducing the overhead of communication and computing

Reducing the overhead of storage

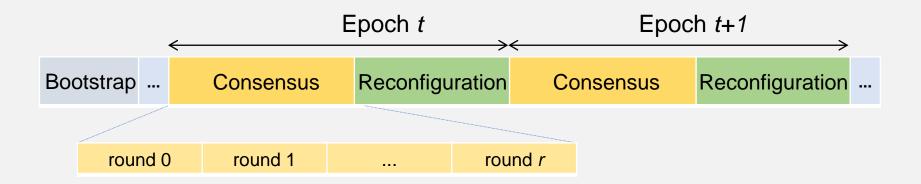


# Motivation



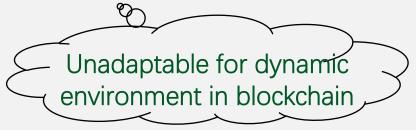


### **1 Privious sharding systems**



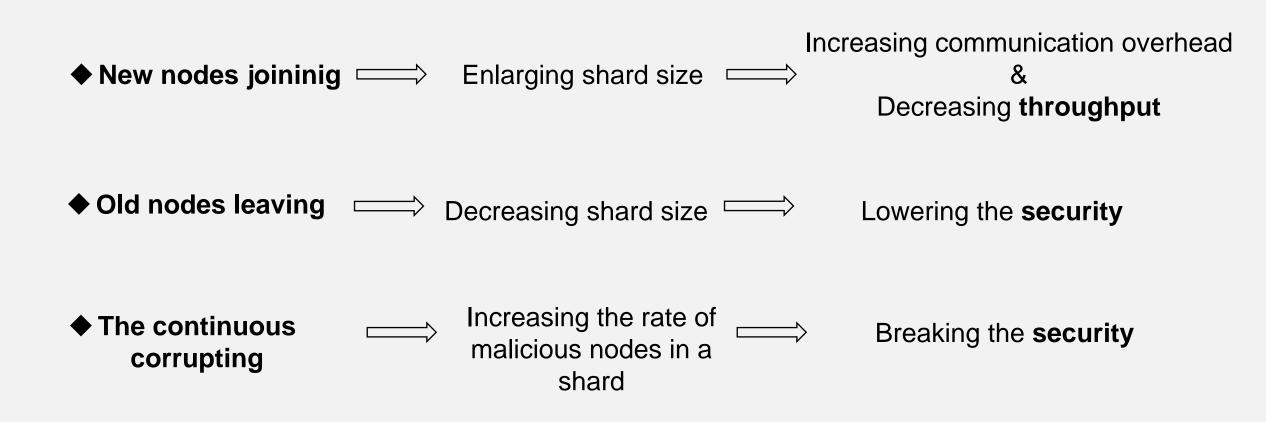
a) Operating in **fixed** interval, called *epoch*. Each epoch includes:

- · Consensus period: consisting of multiple rounds, in which a new block will be created
- Reconfiguration period: reshuffling shards to resist attacks
- b) The number of shards is fixed
- c) The size of block is fixed





### **2** Dynamics of blockchain





# SkyChain A dynamic blockchain sharding system





### 1 SkyChain design

#### **Idea**:

• Adjusting the sharding policy based on the dynamic environment

#### Goal

Achieving a long-term balance between performance and security

#### Challenges

- How to make a suitable sharding policy
- How to efficiently merge or split the ledgers when shard number change



### 2 The sharding policy

#### a) The re-sharding interval

• Determining the duration of consensus period

#### b) The number of shards

• Determining how many shards can process transactions in parallel

#### c) The block size

• Determining how many transactions can be validated in each consensus round



### a) The re-sharding interval

 Less frequent re-sharding can save more time for the consensus process, but it increases the risk of the sharding system.

 More frequent re-sharding can decrease the risk of the sharding system, but it intervenes the consensus process because validators stop processing consensus and suffers from extra cost for the communication and computation of re-sharding operation.



### b) The number of shards

- Less shards can maintain the size of shards and increase the resiliency to malicious attacks.
- More shards can increase the parallelism for processing transactions and reduce communication overhead. However, the probability of forming an unsafe shard will be high if the shard size is small.

#### m: the size of shard

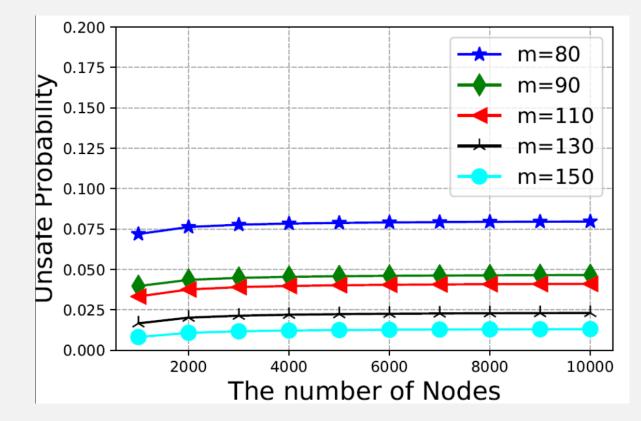


Fig.2:Unsafe probability with shard size and system size



### c) The block size

- A bigger block can pack more transactions, thereby increasing the throughput.
- A bigger block also incurs a significant communication overhead and increases the latency of each consensus round when the shard size is big.

#### m: the size of shard

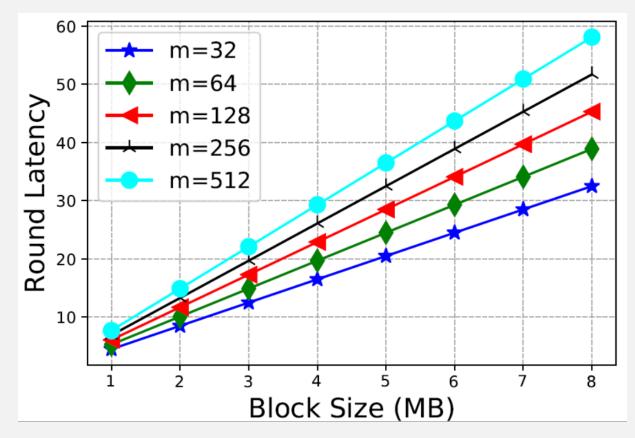


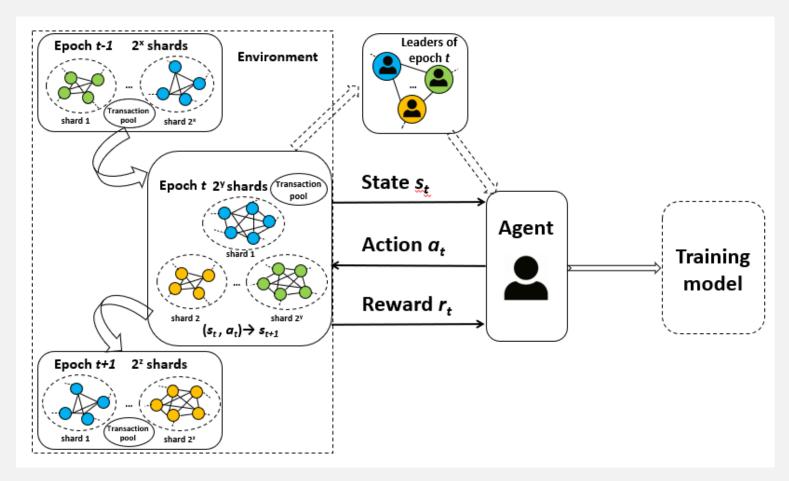
Fig.3:Round latency with block size



## DRL-based dynamic sharding model for making a suitable sharding policy



#### **3 DRL-based sharding model**

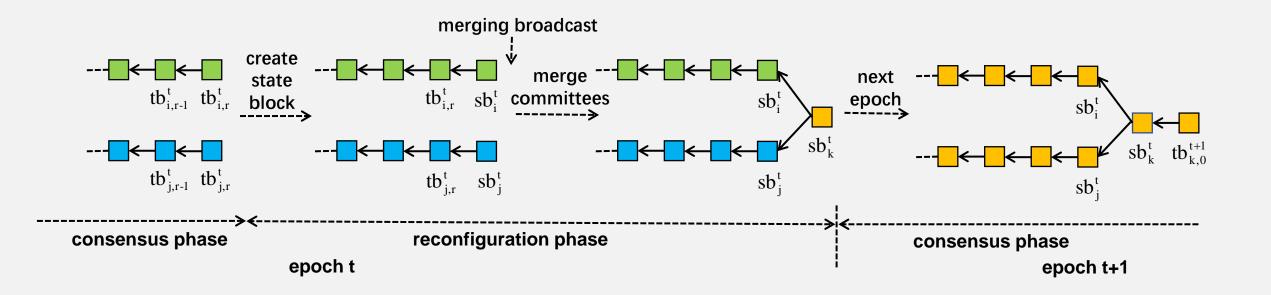




## Adaptive ledger protocol for merging or splitting the ledgers

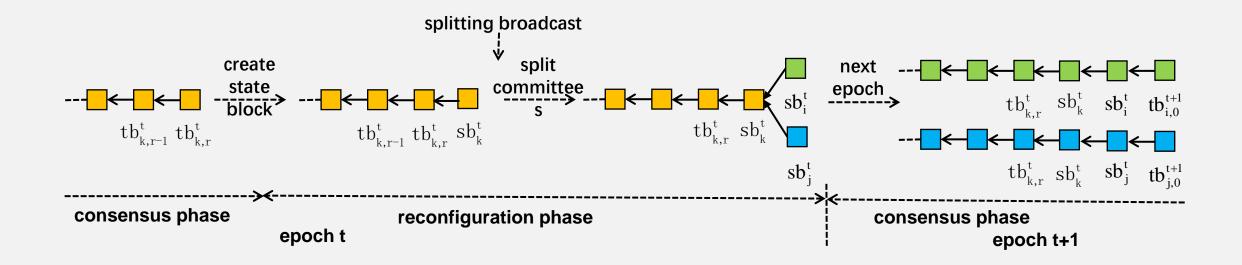


### 4.1 Ledgers merging





### 4.2 Ledgers splitting





# Evaluation





#### **Experimental Setting**

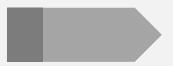
The designs compared sharding algorithms



The sharding policy with fixed re-sharding interval



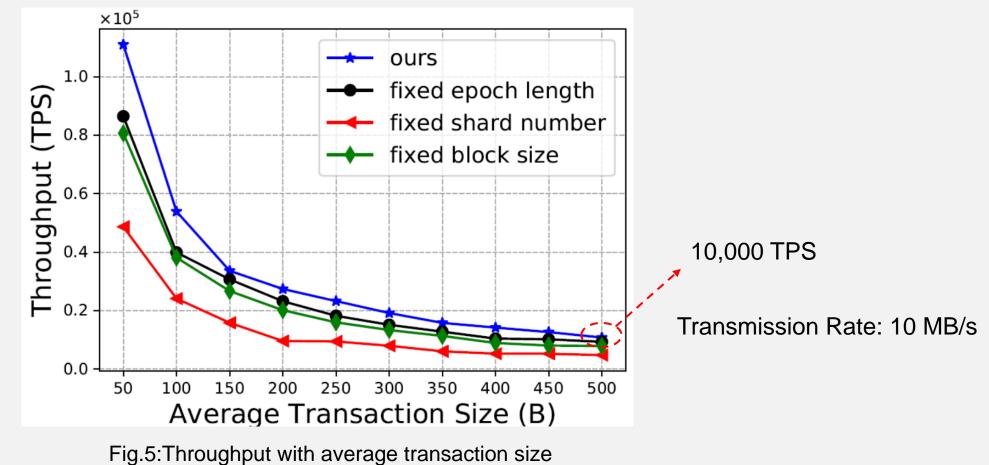
The sharding policy with fixed number of shards



The sharding policy with fixed block size



#### **Throughput vs Transaction size**





#### **Throughput vs Transmission Rate**

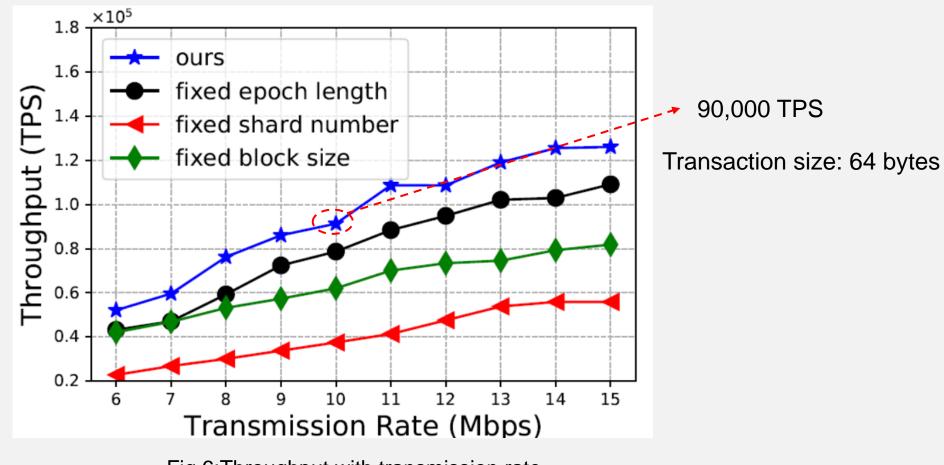


Fig.6:Throughput with transmission rate



## **Thank You!**