

CARD: A Congestion-Aware Request Dispatching Scheme for Replicated Metadata Server Cluster

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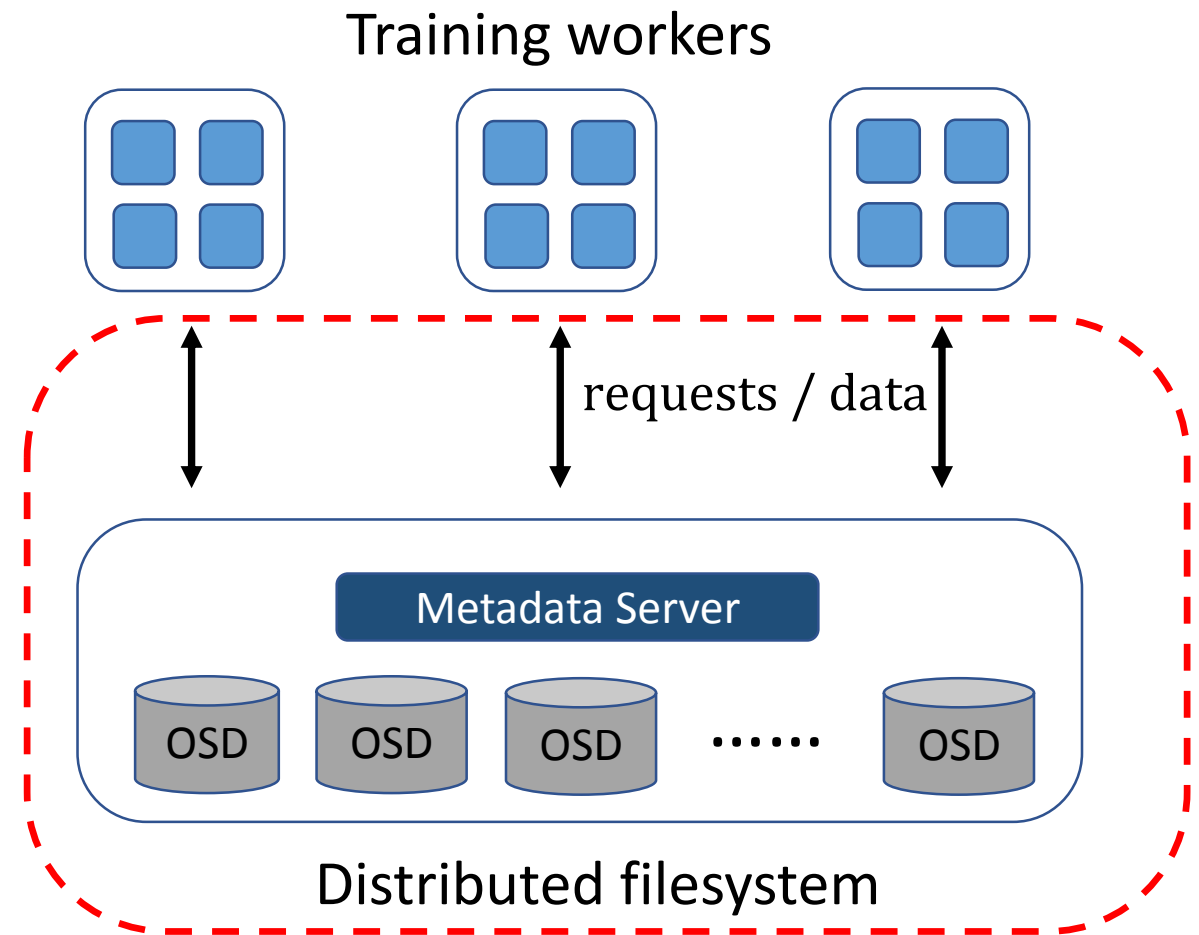


Background: Massive-scale ML in product environments

- Datasets updated hourly or daily
 - data collected and stored in an HDFS-like distributed filesystem
 - periodically offline training for online inference
- Challenges of the data-reader pipeline while training
 - extremely heavy read workloads: millions to billions of files per epoch
 - random access pattern: up-level shuffling for convergence speed

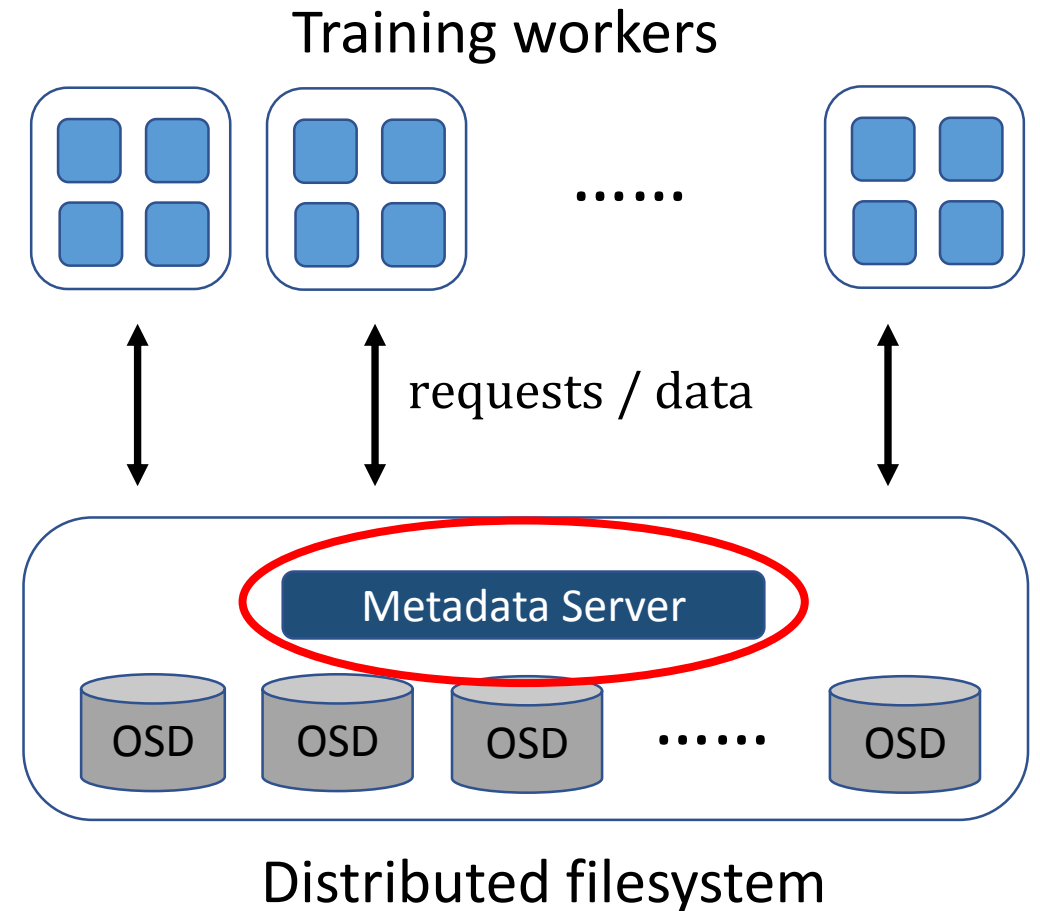
Background: Massive-scale ML in product environments

- Workers interact with a DFS
- Metadata request
-> metadata server (MDS)
- File I/O
-> object storage devices (OSD)



When the number of training workers grows...

- Extremely stressed workloads
- Metadata access step bottlenecks the data-reader pipeline
- Potential single point of failure on MDS



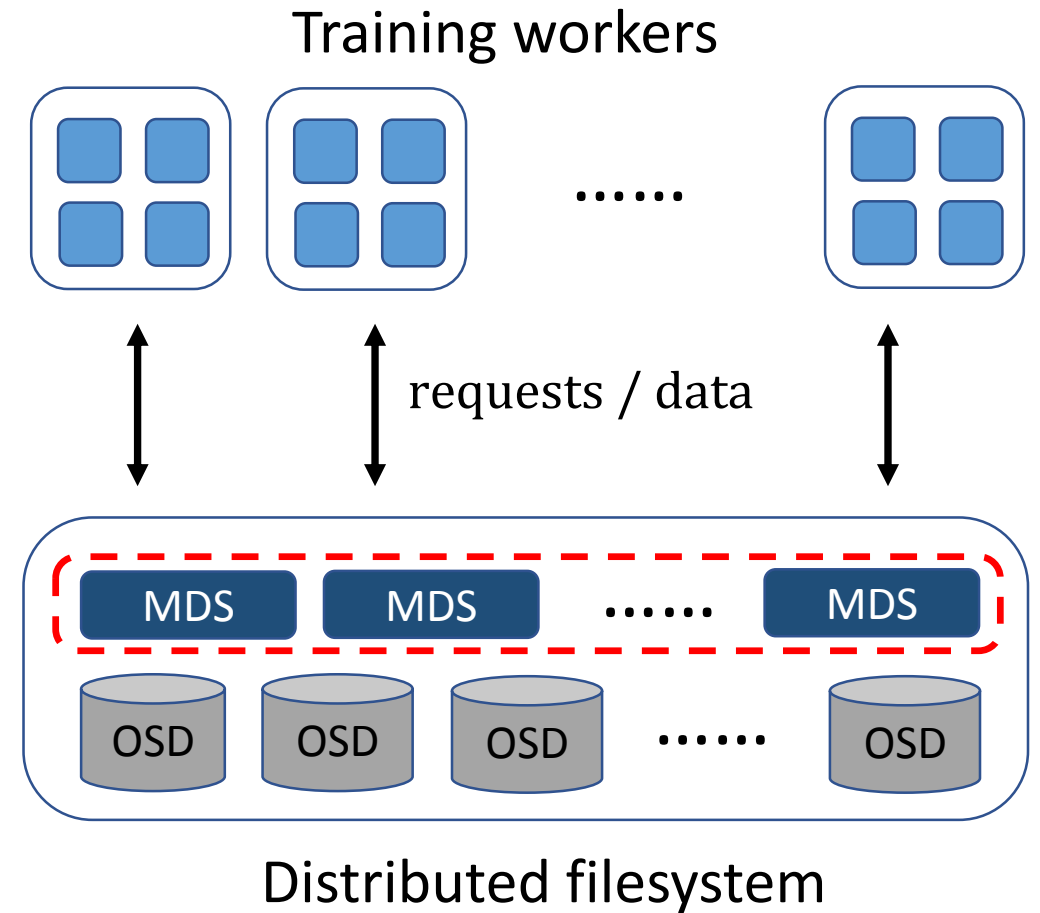
Typical industrial response: Scaling out likewise

- Concerns to be addressed:

- Cost-effectiveness

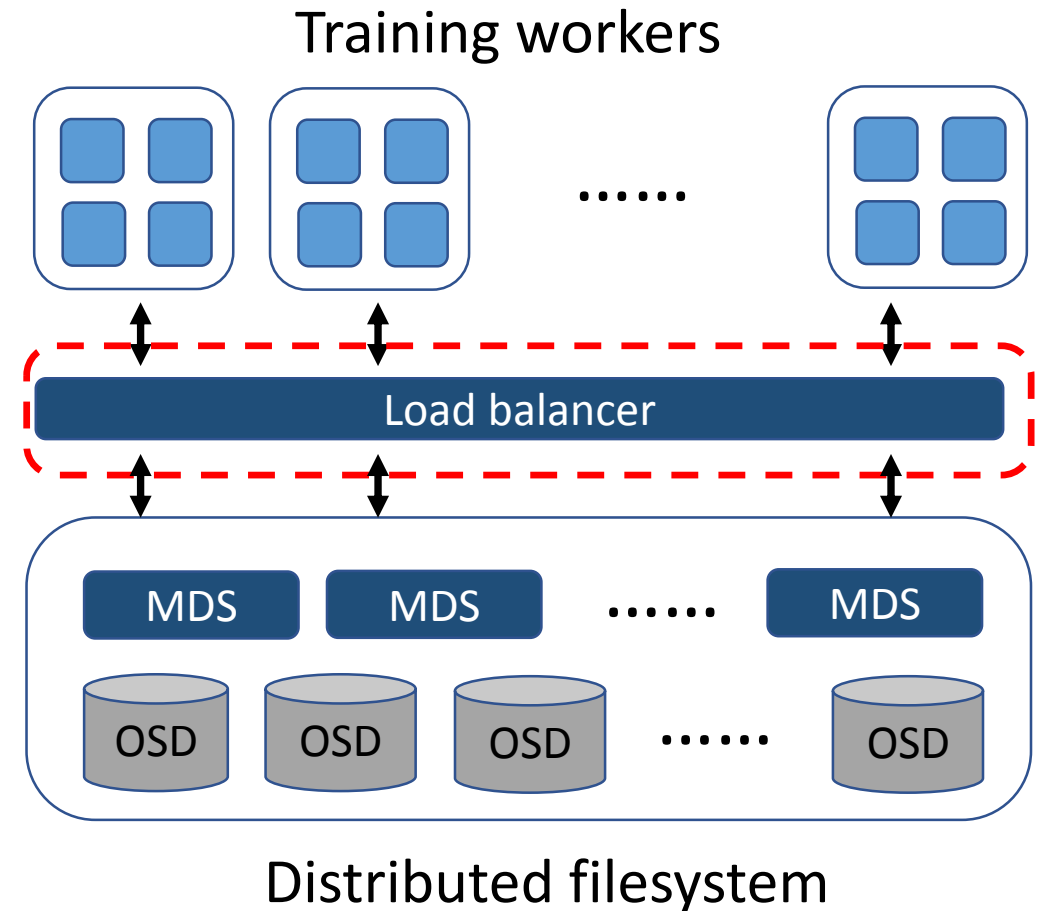
- Scalability

- Run-time stability



To achieve load-balance...

- A middle layer load-balancer
 - Pros:
 - good global load balancing
 - more features are optional
 - Cons:
 - load-balancer is stressed
 - reintroduce a potential single point of failure
 - not cost-effective



To achieve load-balance...

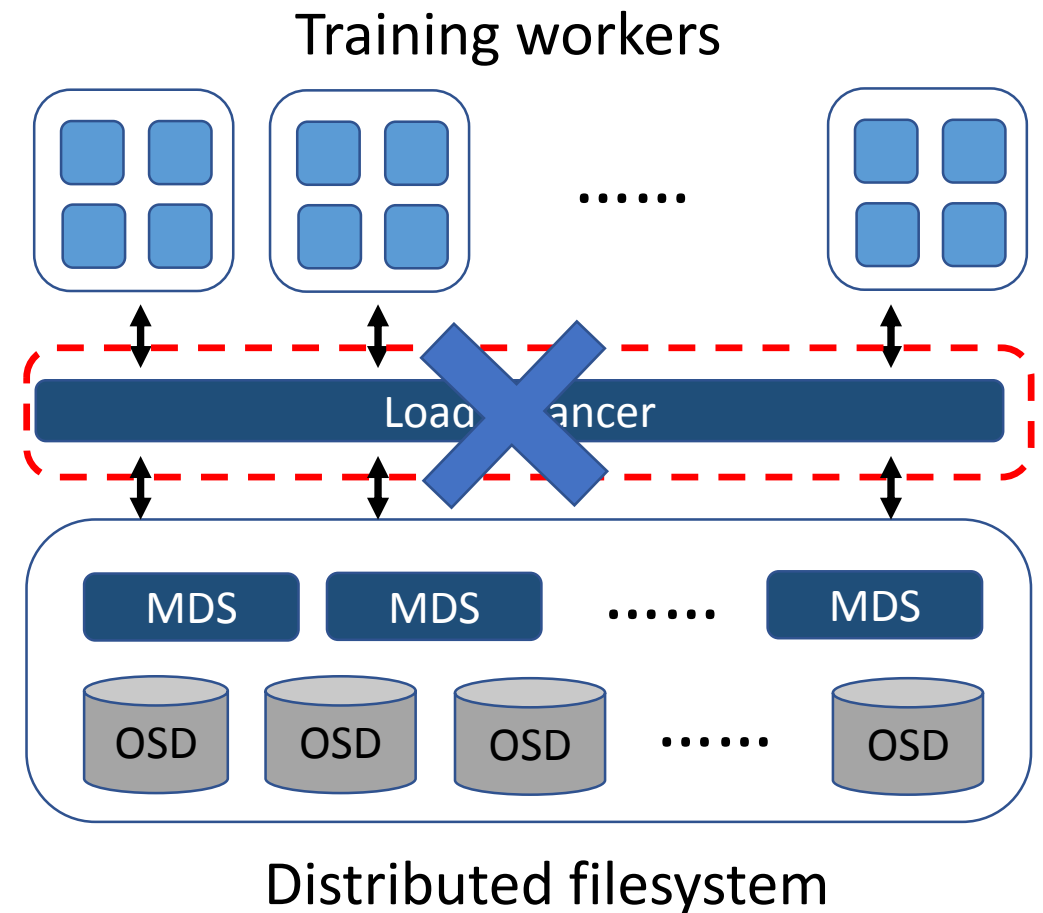
- A middle layer load-balancer

- Pros:

- good global load balancing
- more features are optional

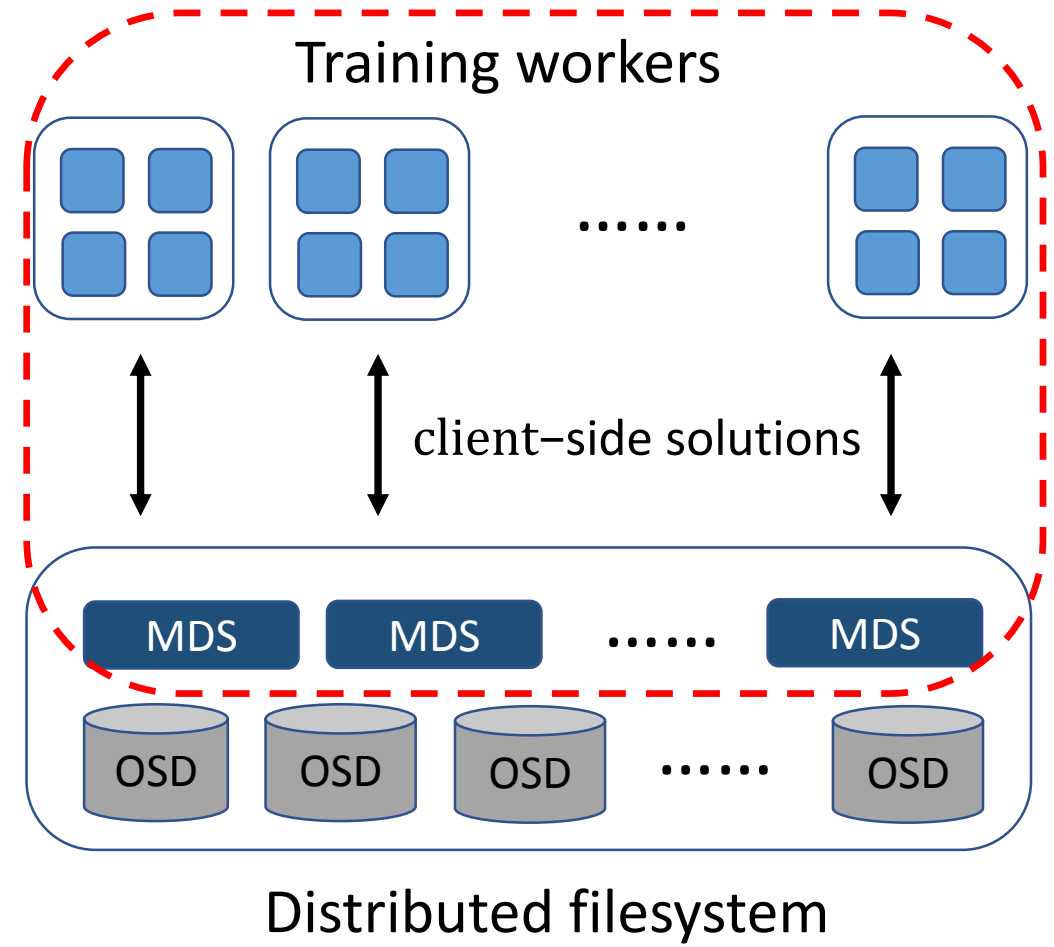
- Cons:

- load-balancer is stressed
- reintroduce a potential single point of failure
- not cost-effective



Try client-side solutions

- Easy to implement
- Cost-effective



Client-side solution: Round-Robin

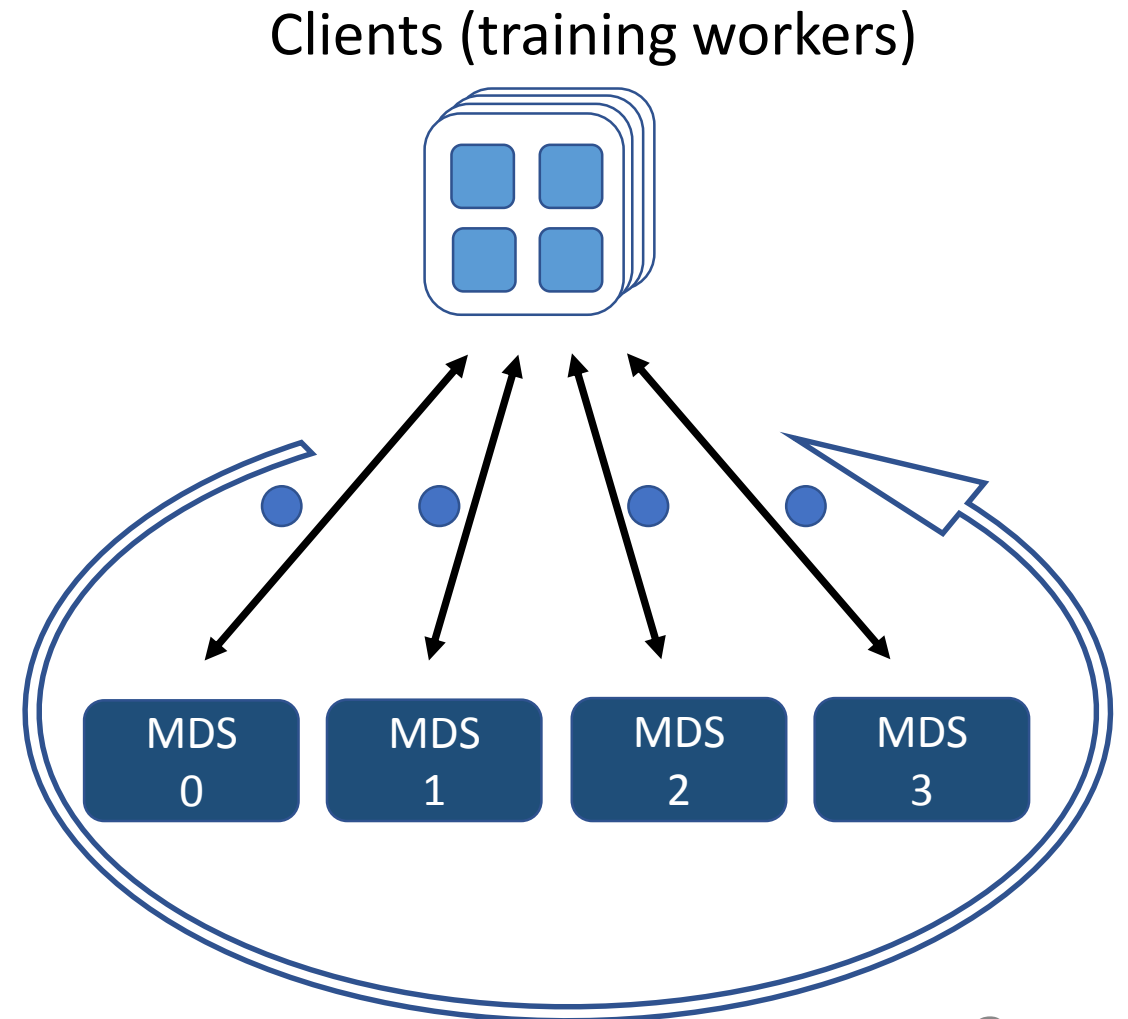
- Round-Robin

- Pros:

- simple yet effective in homogeneous environments

- Cons:

- inflexible and inefficient in shifting or heterogeneous environments



Client-side solution: Heuristic selection

- Heuristic selection

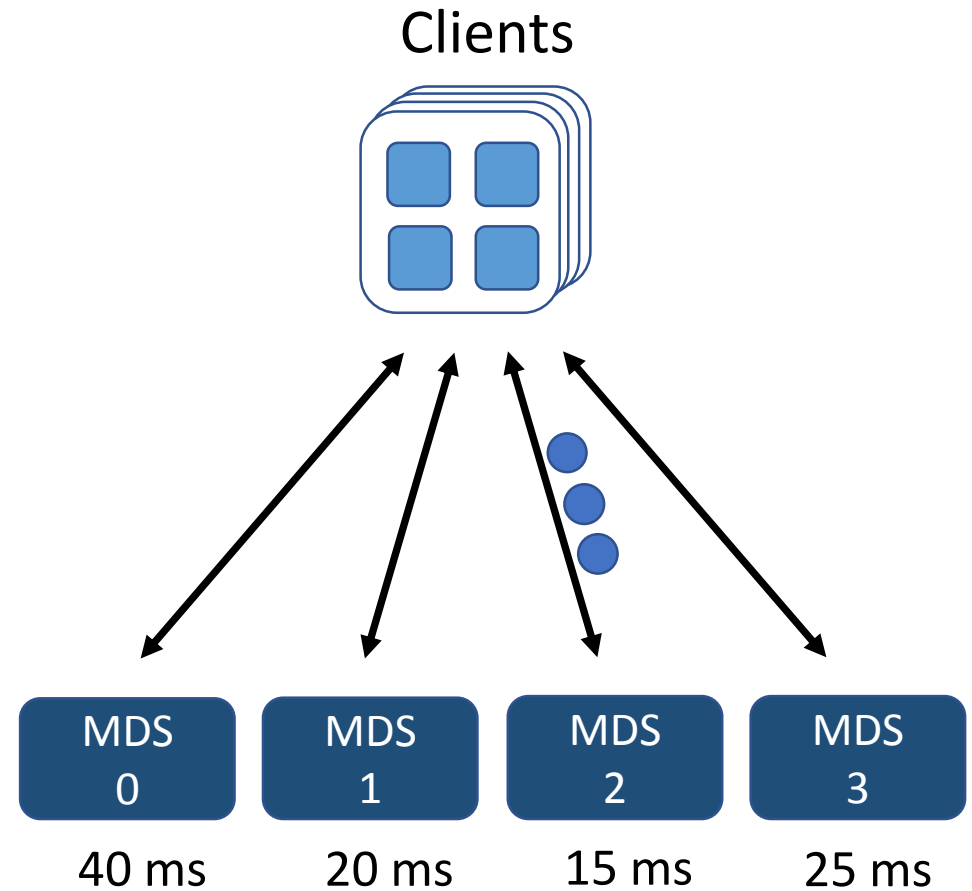
- e.g., prefer lowest MART (moving average of response time)

- Pros:

- effective when facing light-weight workloads

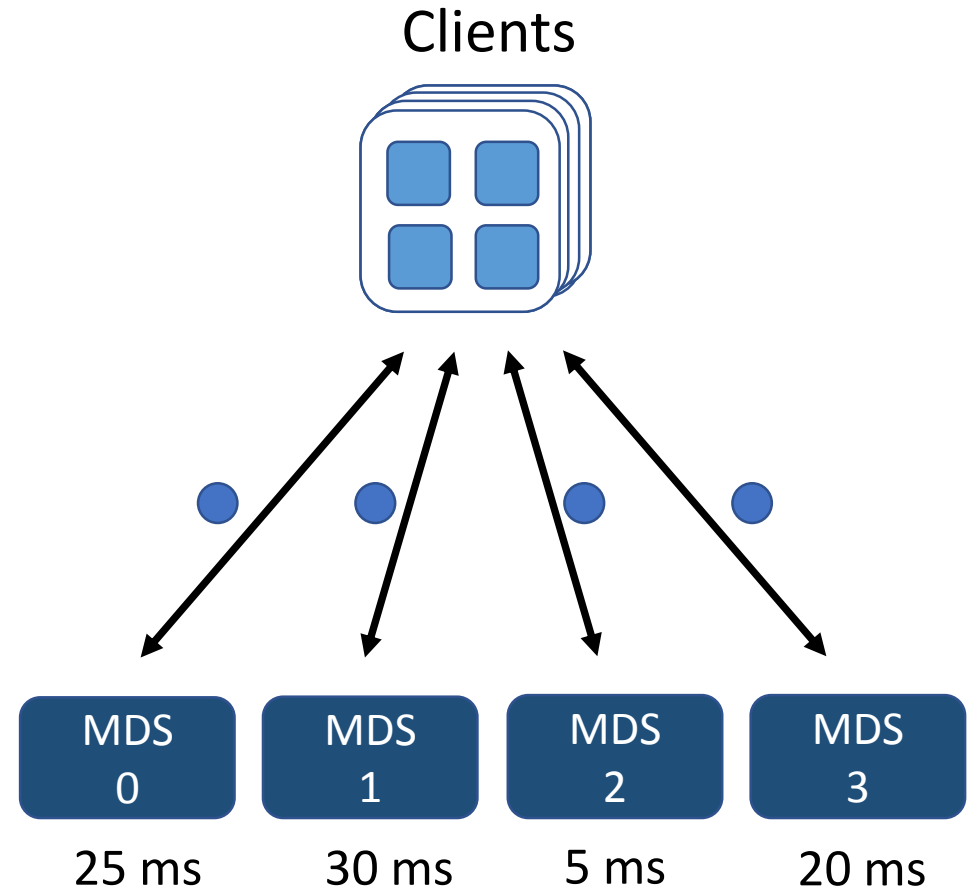
- Cons:

- cause herd-behavior and load-oscillations



Client-side solution: Round-Robin with Throttling

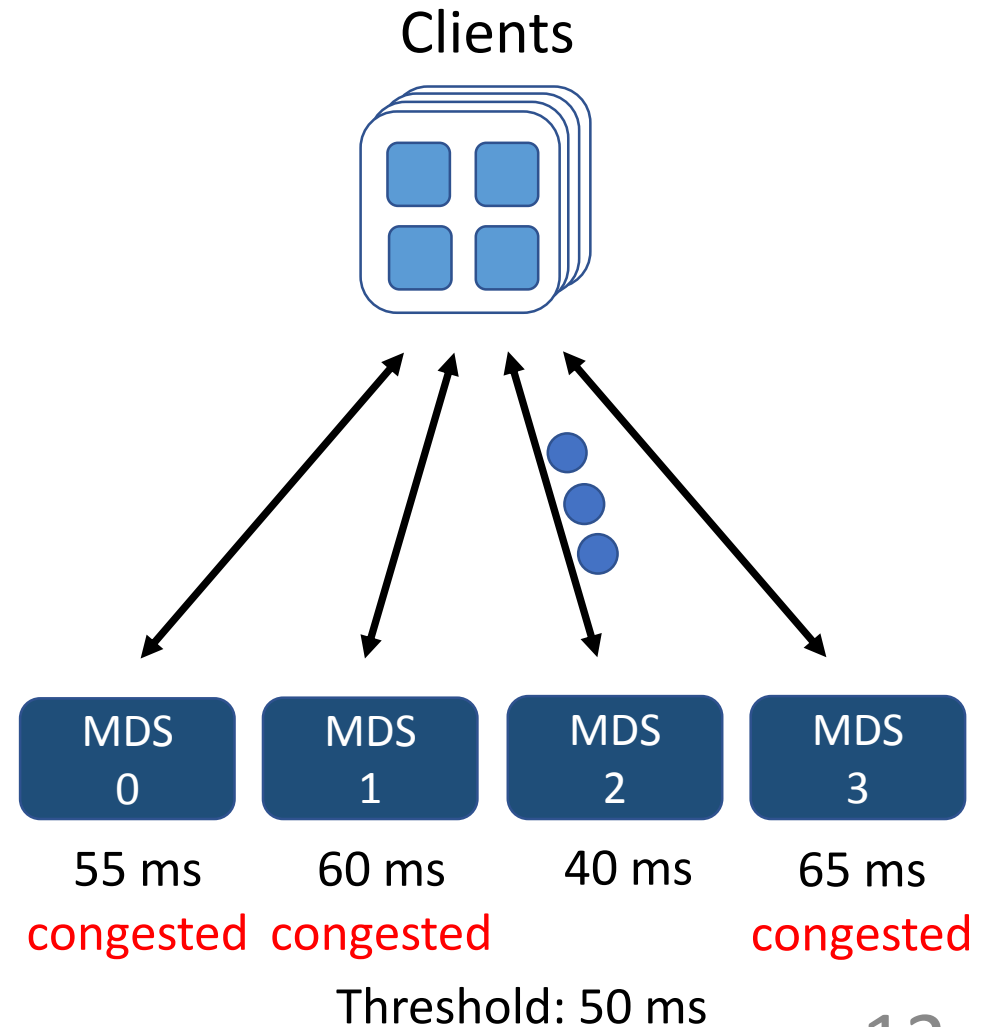
- Round-Robin with throttling
 - e.g., LADS, preset a MART threshold to mark servers as congested
- Light-weight workloads
 - = Round-Robin



Threshold: 50 ms

Client-side solution: Round-Robin with Throttling

- Round-Robin with throttling
 - e.g., LADS, preset a MART threshold to mark servers as congested
- Light-weight workloads
 - = Round-Robin
- Heavy workloads
 - = Heuristic selection
 - herd-behavior and load-oscillations remain

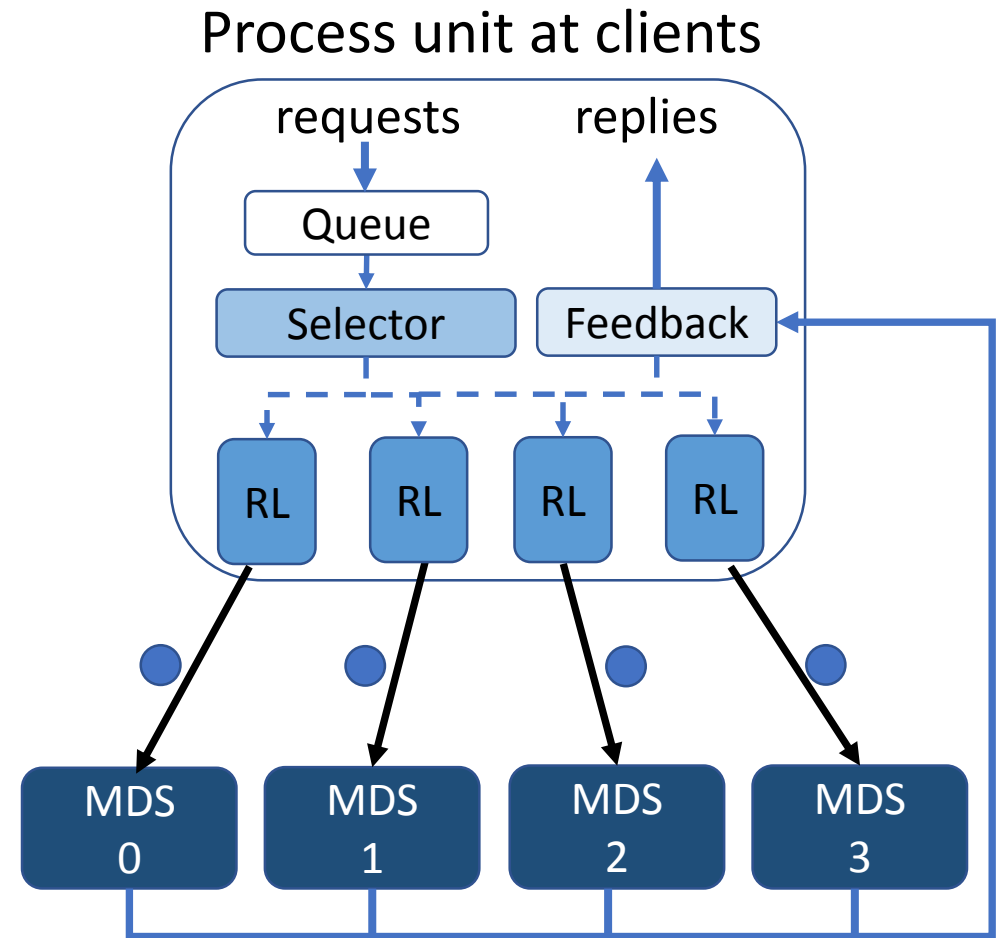


CARD: Congestion-Aware Request Dispatching scheme

- Core idea: Round-Robin with adaptive rate-control
 - inspired by CUBIC for TCP protocol
 - counting-based implementation
 - no extra info required from servers
- Light-weight workloads
 - = Round-Robin
- Heavy workloads
 - redirect requests from overloaded MDS to underloaded MDS
 - suppress upcoming requests: if and only if all servers are overloaded

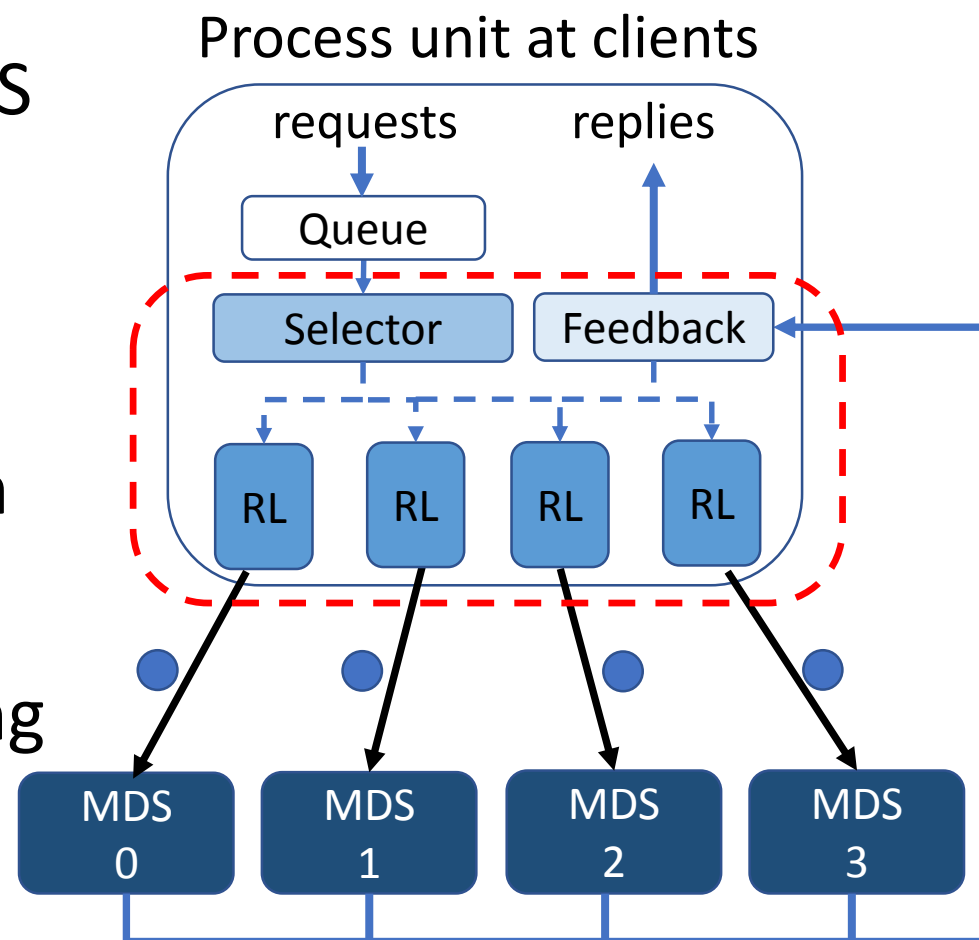
Congestion-aware rate-control mechanism

- Queue: place pending requests
- Selector: Round-Robin dispatching
- Rate-limiter: rate-control module
- Feedback: process feedbacks and forward replies



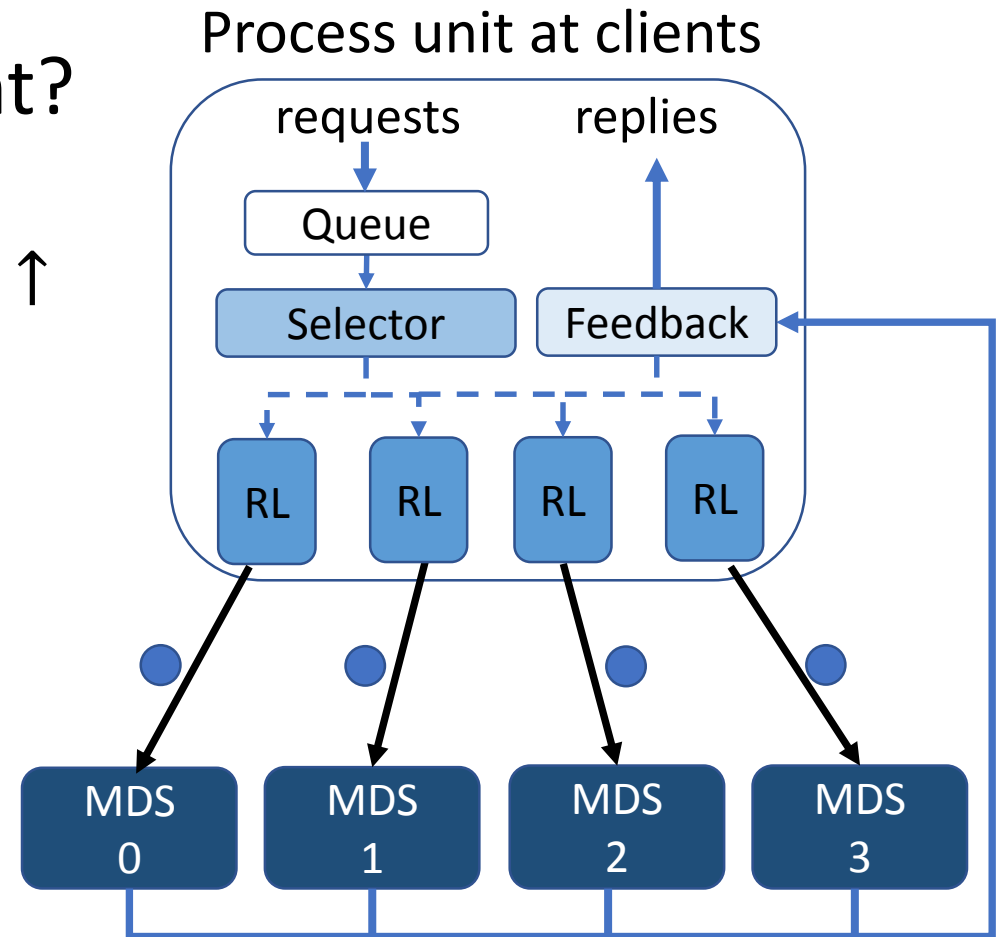
Congestion-aware rate-control mechanism

- Restrict requests routed to each MDS per δ time window
- Gradually increase the restriction according to a cubic growth function
- Feedback module computes receiving rates after each time window and forwards to RLs



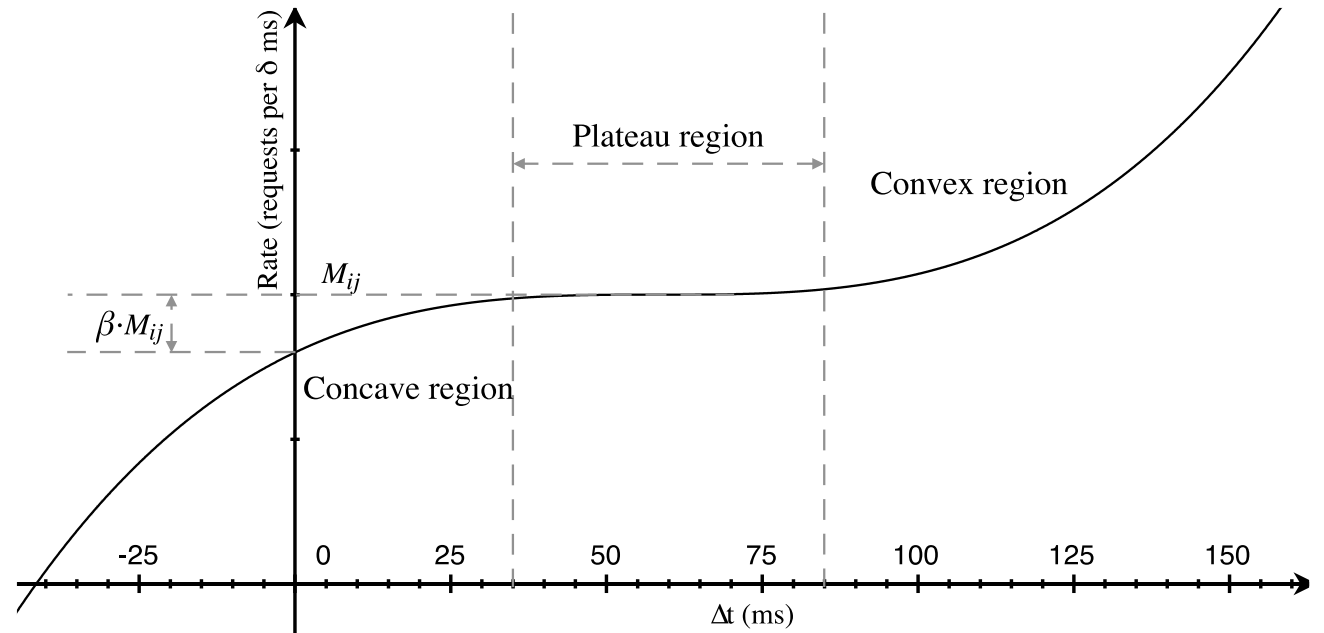
Congestion-aware rate-control mechanism

- How to identify a congestion event?
 - sending rate $>$ receiving rate
 - elapsed time since last sending rate \uparrow event $> \lambda$ (a hysteresis period)
- What to do then?
 - record current sending rate as saturated sending rate
 - reduce current sending rate



The cubic growth function for the rate-control

- Δt : elapsed time since the last congestion event
- M_{ij} : saturated sending rate
 - Changed to current sending rate adaptively whenever a congestion event happens
 - Then, current sending rate reduced to $(1 - \beta) \cdot M_{ij}$, and start to grow all over again accordingly



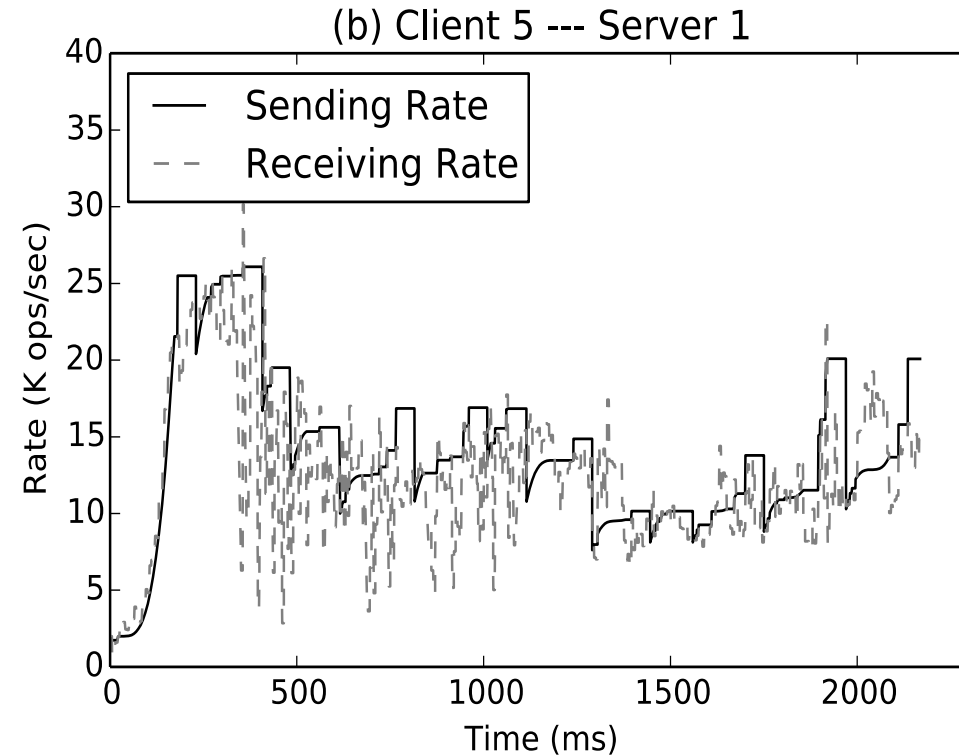
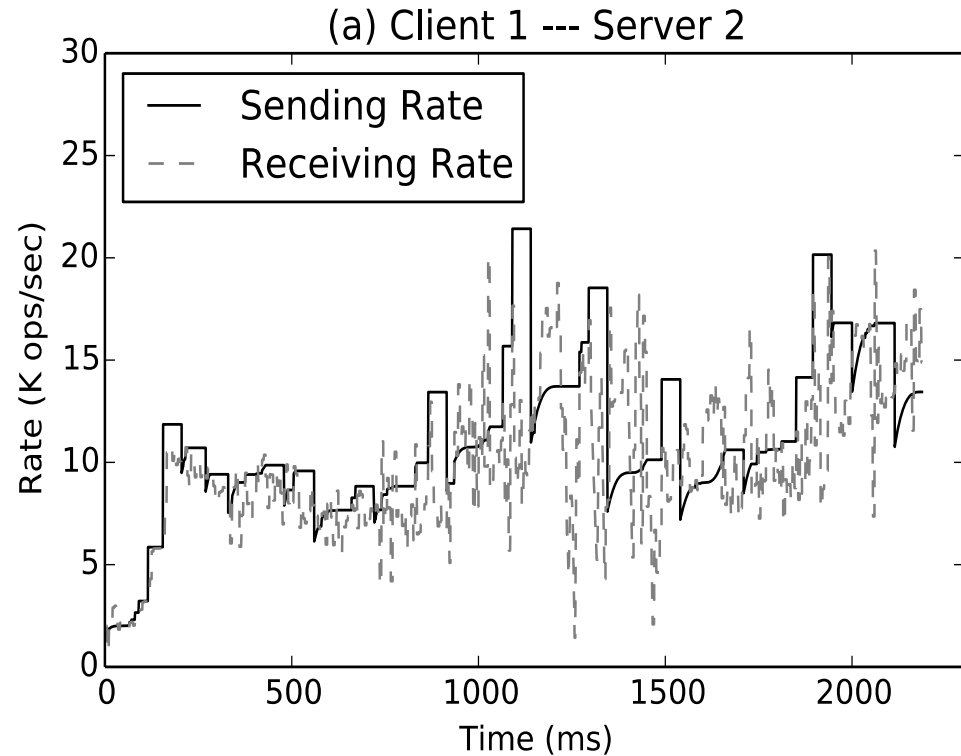
Evaluation setup

- We implemented a prototype RMSC for simulation purposes
- Up to 8 servers to measure system scalability
- Crafted descending setup for heterogeneous experiments
- 10 clients run on separate machines launching request with Poisson arrivals
- $\delta = 5 \text{ ms}$, $\lambda = 10 \text{ ms}$, $\beta = 0.20$
- To compare against CARD, we implemented aforementioned Round-Robin, MART and LADS as well
- Refer to the paper for more setup details

Evaluation highlights

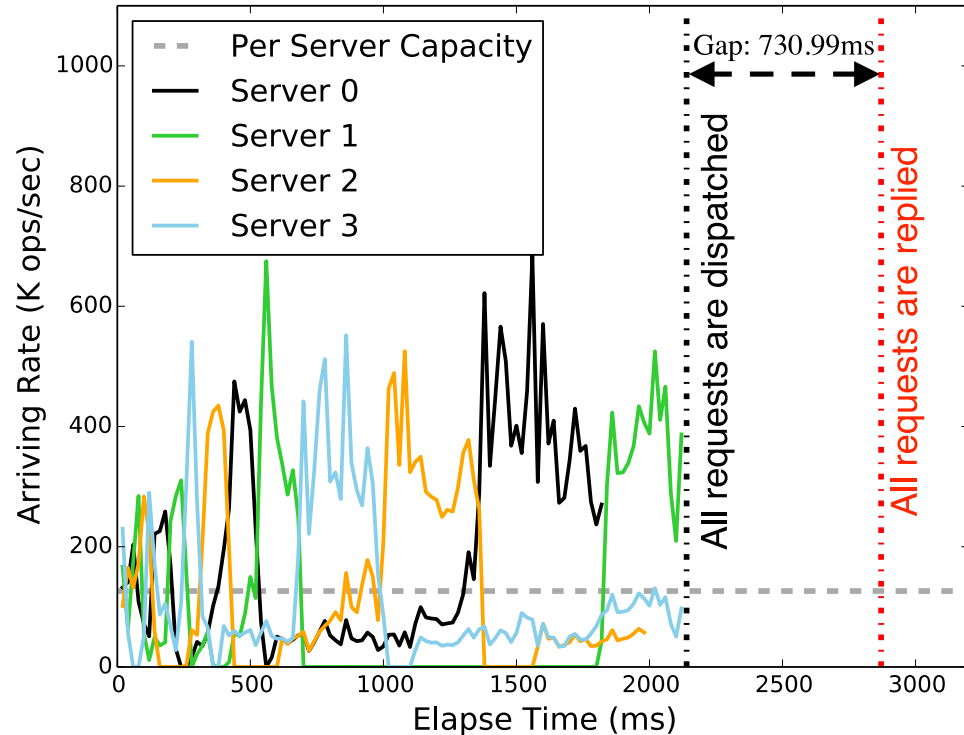
- Do CARD's rate-control mechanism work as expected?
 - Yes, the rate-control process is effective and adaptive
 - Loads among servers are balanced under heavy workloads
- Can CARD achieve better scalability?
 - In homogeneous clusters: $\text{CARD} \approx \text{Round-Robin} > \text{other strategies}$
 - In heterogeneous clusters: Yes, $\text{CARD} > \text{other strategies}$

Examples of the rate-control procedure

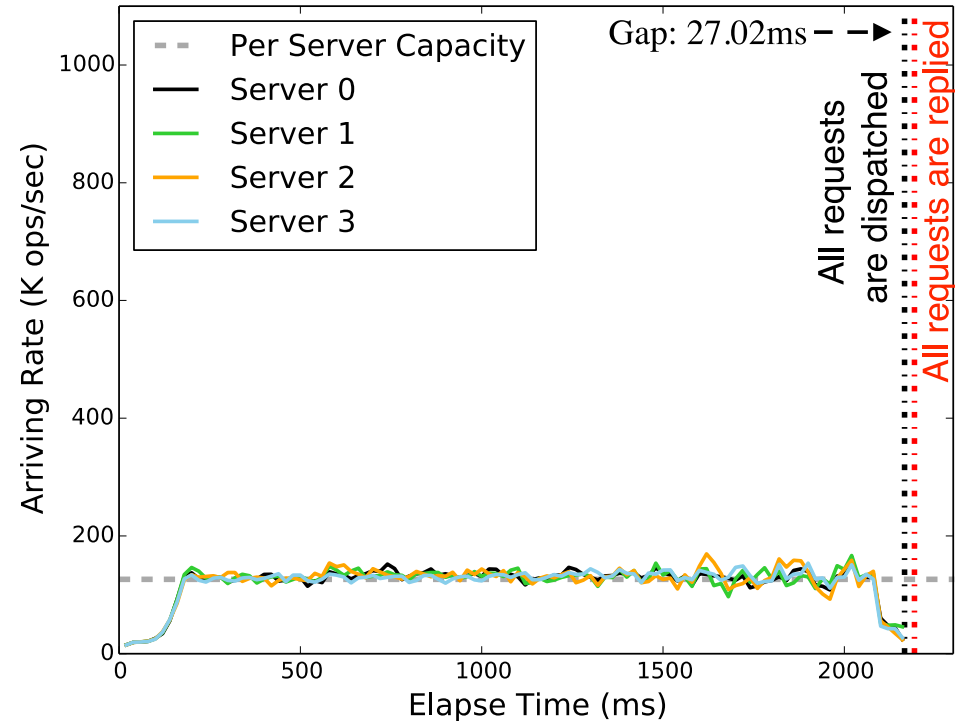


The sending rate from each client to each server is adjusted adaptively according to the receiving rate

Overall arriving rates in the homogeneous cluster



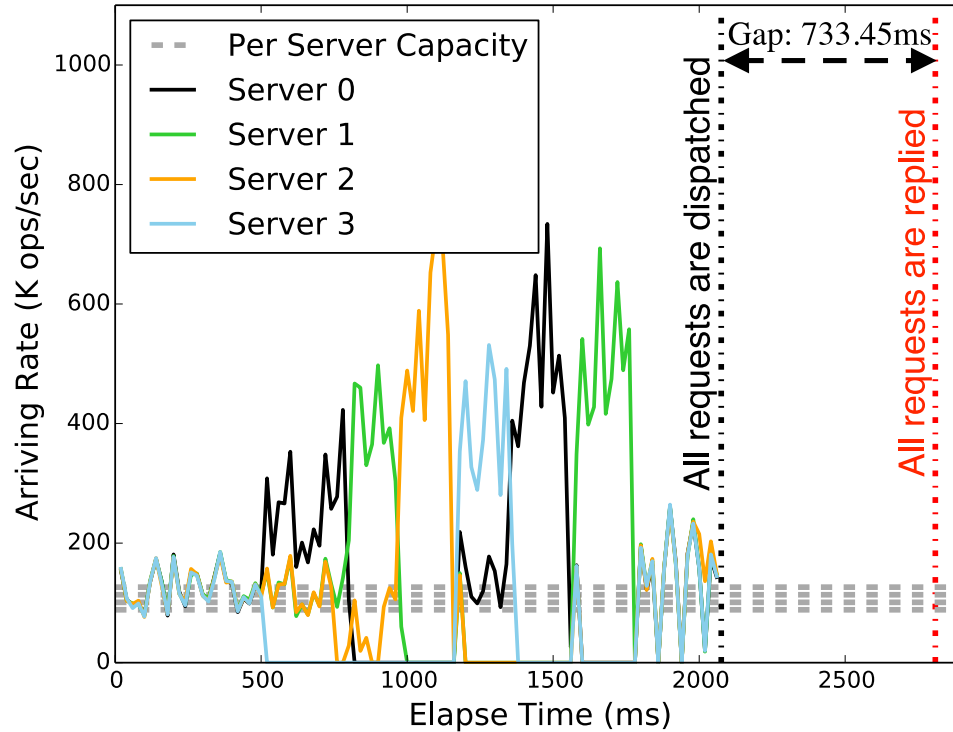
MART



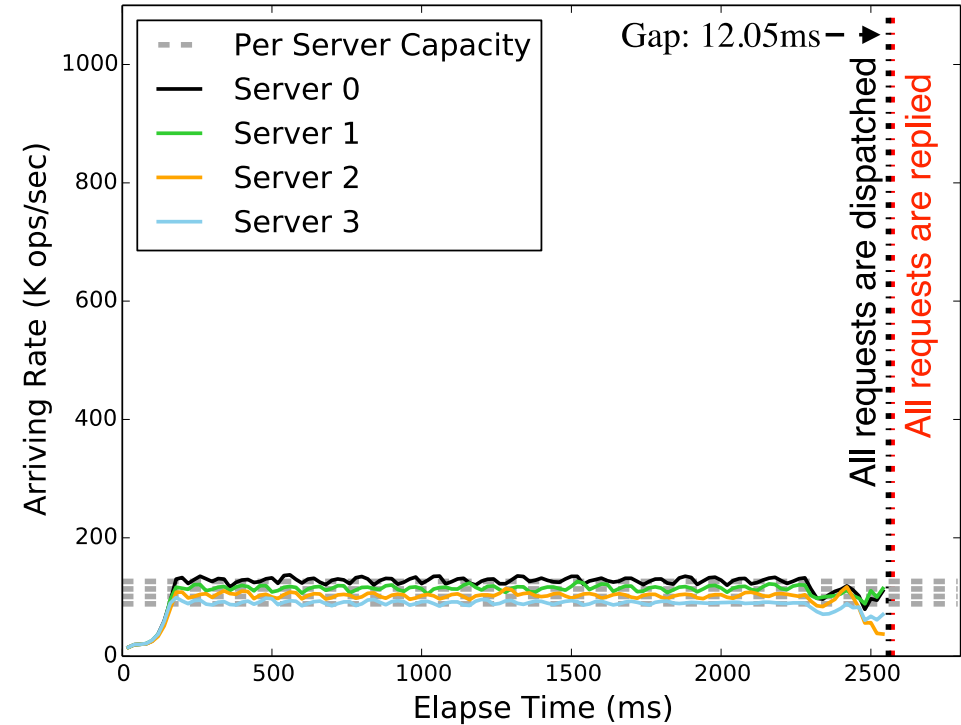
CARD

- 1) Heuristic selections cause severe herd behavior and load-oscillations
- 2) A data loading job is completed earlier when using CARD

Overall arriving rates in the heterogeneous cluster



LADS

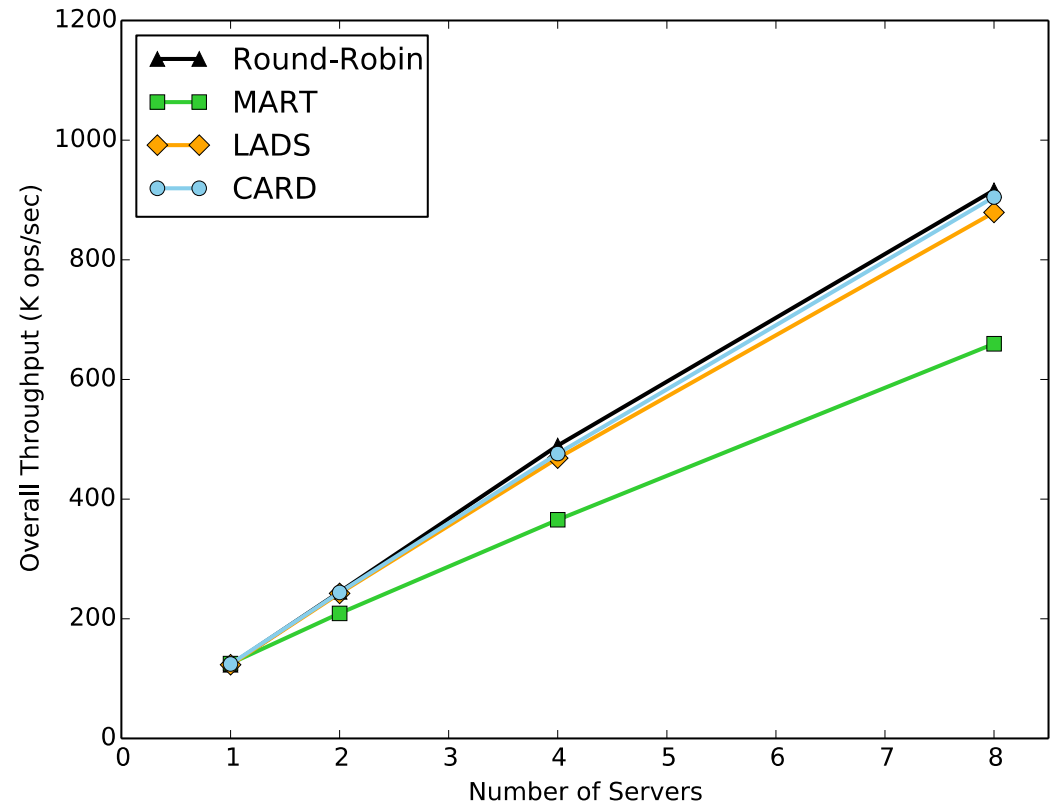


CARD

- 1) A basic threshold throttling strategy is not sufficient enough
- 2) Arriving rates are stabilized around servers' capacity when using CARD

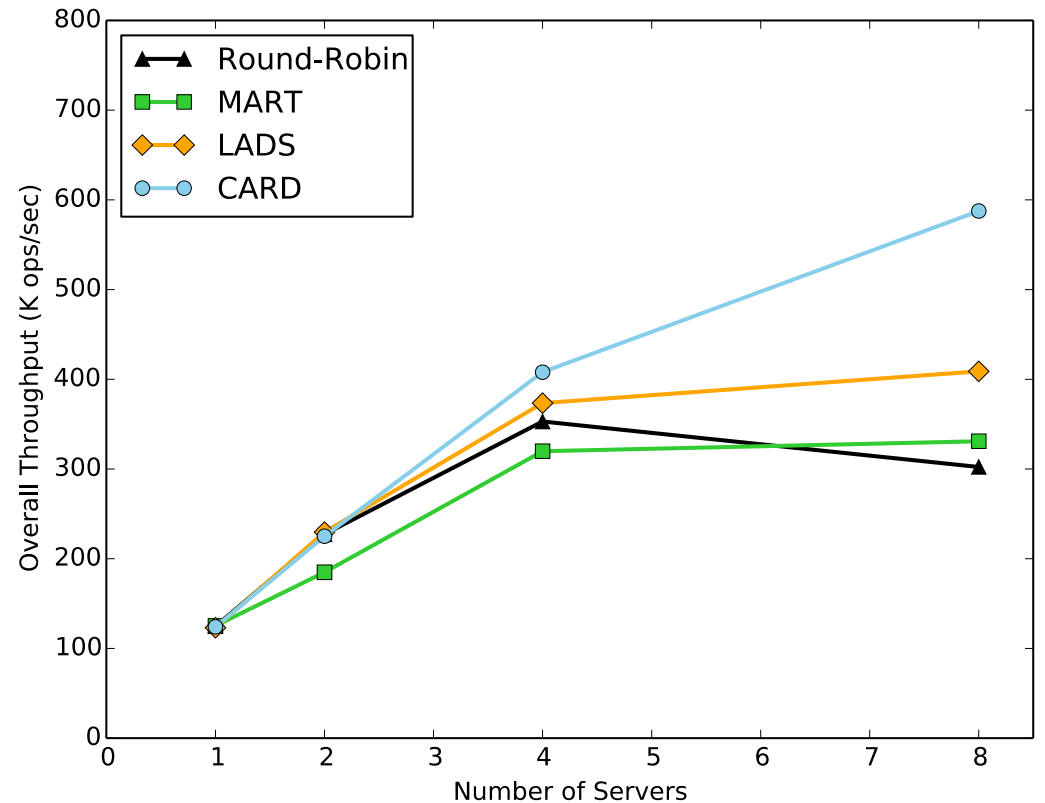
Overall throughput in the homogeneous cluster

- Heuristic selection is a bad choice under heavy workloads
- In ideal homogenous environments, Round-Robin and CARD achieve great scalability



Overall throughput in the heterogeneous cluster

- Round-Robin is ineligible when facing heterogenous setups
- CARD outperforms other strategies and achieves excellent scalability



Summary: CARD

- Adaptive client-side throttling method: easy and efficient
- Redirect requests from the overloaded server to the underloaded server adaptively under heavy workloads
- Degrade into pure Round-Robin when facing light-weight workloads
- Boosts throughput significantly over competing strategies in heterogeneous environments